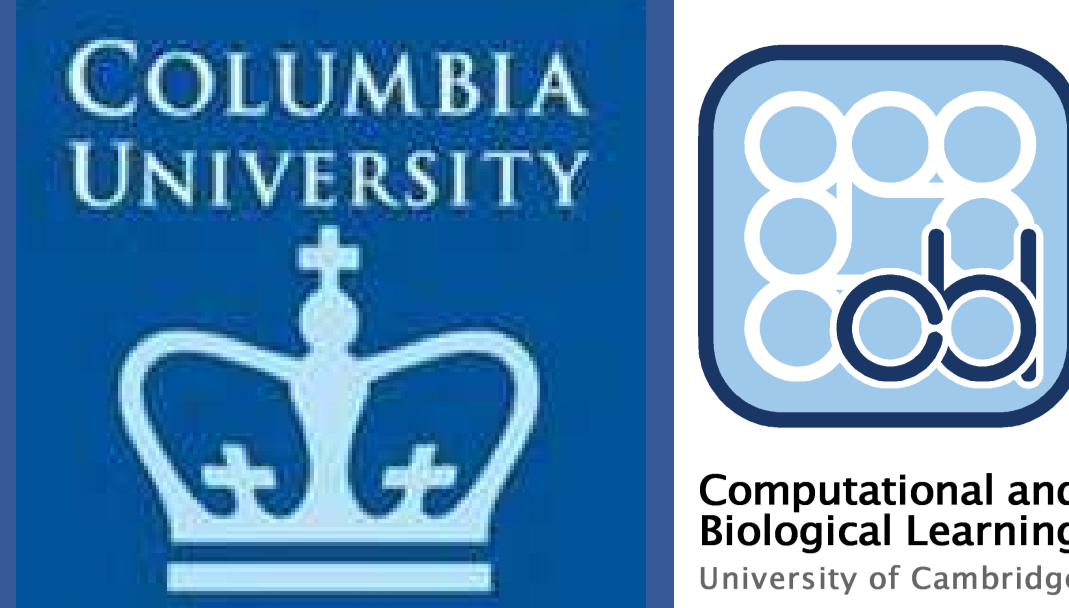




Clamping Variables and Approximate Inference

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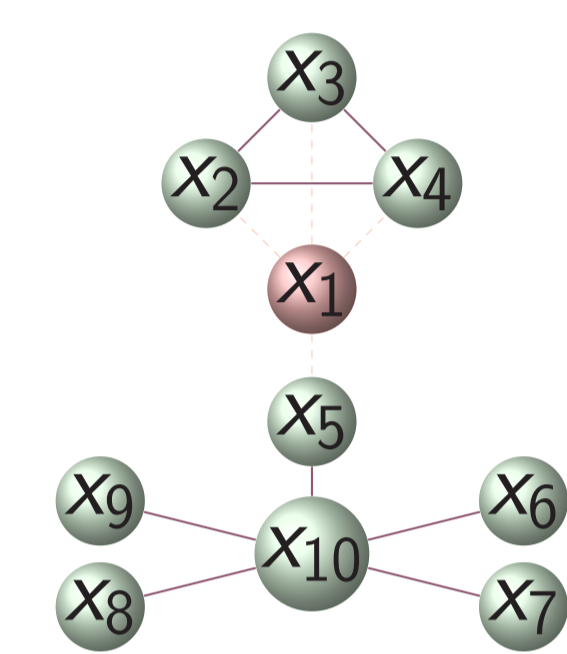


Summary

We address the problem of *marginal inference* for undirected graphical models - estimating the partition function Z and marginal probability distributions. We focus on *binary pairwise* (Ising) models, e.g. vision, RBMs, or social networks. Combining *clamping* of variables with *approximate variational inference* we obtain

- Strong theoretical results (middle column)
- Promising empirical results (right column)

Background: What is clamping?



Z can be split into two parts: **clamp** variable X_1 to each of $\{0, 1\}$, then add the two sub-partition functions:

$$Z = Z|_{X_1=0} + Z|_{X_1=1}$$

After clamping a variable, remove it from the graph

- If remaining sub-models are **acyclic** then can find sub-partition functions efficiently (Bethe approximation is exact on trees)

- If not,
 - (a) Can repeat until acyclic, *or*
 - (b) Settle for **approximate inference** on sub-models

Will clamping and summing approximate sub-partition functions always lead to a better estimate of Z than approximate inference on the original model?

Often but not always (see paper for example)

$x_1 x_2 \dots x_{10}$	score	$\exp(\text{score})$
0 0 ... 0	1	2.7
0 0 ... 1	2	7.4
...
0 1 ... 1	1.3	3.7
1 0 ... 0	-1	0.4
1 0 ... 1	0.2	1.2
...
1 1 ... 1	1.8	6.0
Total $Z =$	47.1	

Variational inference

$$p(x) = \frac{1}{Z} \exp(\theta \cdot x)$$

- Exact inference may be viewed as *optimization*,

$$\log Z = \max_{\mu \in \mathcal{M}} [\theta \cdot \mu + S(\mu)], \quad S \text{ is true entropy}$$

- Bethe makes 2 pairwise approximations,

$$\log Z_B = \max_{q \in \mathcal{L}} [\theta \cdot q + S_B(q)], \quad S_B \text{ is Bethe entropy}$$

- Bethe is exact on trees

- Observe that when X_i is clamped, we optimize over a subset

$$\log Z_B|_{X_i=0} = \max_{q \in \mathcal{L}: q_i=0} [\theta \cdot q + S_B(q)], \quad q_i = q(X_i = 1)$$

$$\Rightarrow Z_B|_{X_i=0} \leq Z_B, \text{ similarly } Z_B|_{X_i=1} \leq Z_B$$

- **Notation:** If we *clamp* variable X_i and *sum approximate* sub-partition functions,

$$Z_B^{(i)} := Z_B|_{X_i=0} + Z_B|_{X_i=1} \leq 2Z_B \text{ by above}$$

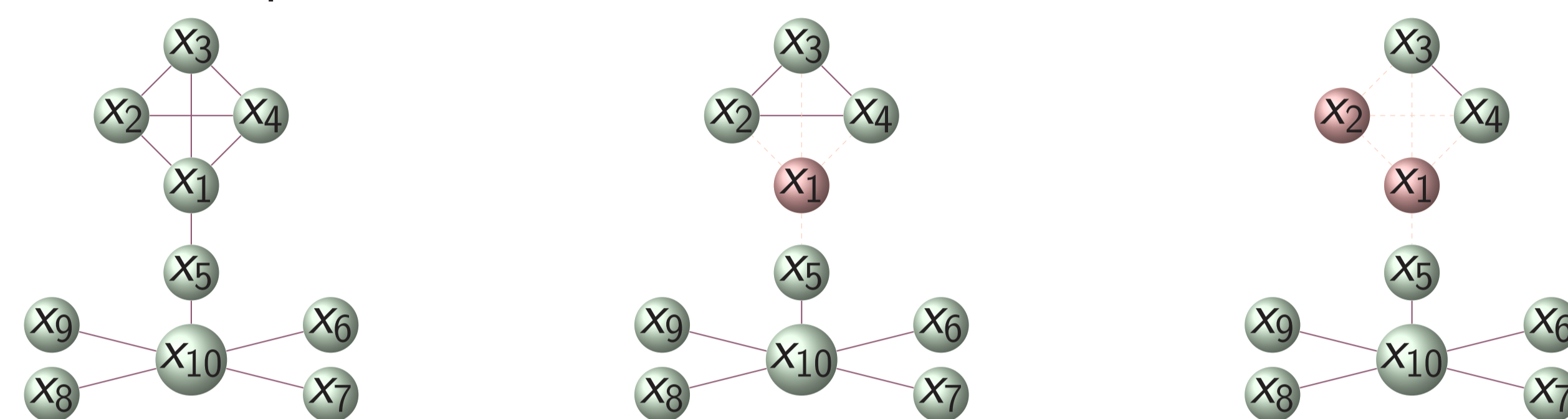
General models: upper bound on Z

Upper bound on $Z_B^{(i)}$ leads to upper bound on Z

- $Z_B^{(i)} := Z_B|_{X_i=0} + Z_B|_{X_i=1} \leq 2Z_B$
- Repeat, clamping variables until remaining model is acyclic, where Bethe is exact
- For example, if we must delete 2 variables X_i, X_j , obtain

$$Z_B^{(ij)} := \sum_{a,b \in \{0,1\}} Z_B|_{X_i=a, X_j=b} \leq 2^2 Z_B$$

But sub-partition functions are *exact*, hence $\text{LHS} = Z$



Let $\nu(G)$ be minimum size of a **feedback vertex set** (set of vertices such that deleting them renders graph acyclic; $\nu \geq \text{treewidth} - 1$)

Theorem (result is tight)

$$Z \leq 2^\nu Z_B$$

Attractive models: Master Theorem (strongest result)

An *attractive* model has all variables positively coupled

- For any variable X_i and $x \in [0, 1]$, let $q_i = q(X_i = 1)$ and

$$\log Z_{B_i}(x) = \max_{q \in \mathcal{L}: q_i=x} [\theta \cdot q + S_B(q)]$$

- $Z_{B_i}(x)$ is 'Bethe partition function constrained to singleton $q_i = x$ '
- Define new function,

$$A_i(q_i) := \log Z_{B_i}(q_i) - S_i(q_i)$$

- By considering derivatives of the Bethe free energy, and how the optimum constrained to singleton q_i varies with q_i , we show

Theorem (strongest result for attractive models)

For an attractive binary pairwise model, $A_i(q_i)$ is convex

Attractive models: Consequences of Master Theorem

Lower bound on $Z_B^{(i)}$ leads to lower bound on Z

Theorem (clamp and sum can only increase Bethe)

For an attractive binary pairwise model and any X_i , $Z_B \leq Z_B^{(i)}$

Then with similar proof to result above for general models,

Corollary (lower bound on Z , first proved by Ruozi, 2012)

For an attractive binary pairwise model, $Z_B \leq Z$

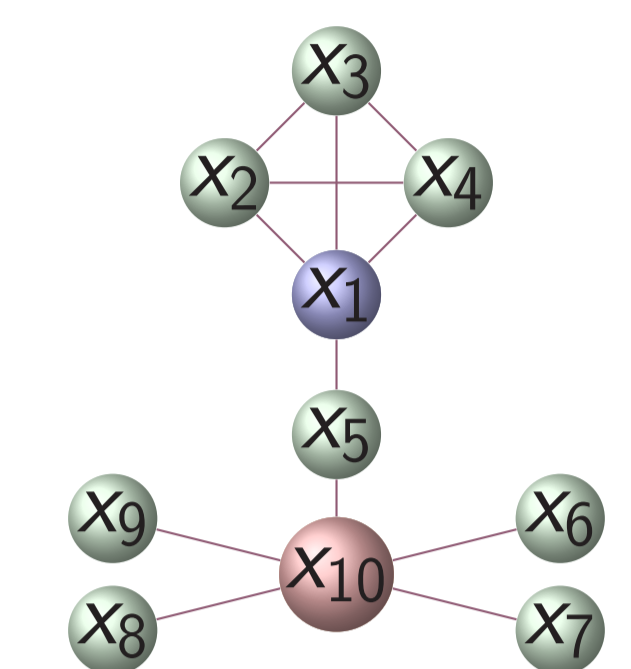
\Rightarrow clamping can only *improve* the estimate of the partition function

$$Z_B \leq Z_B^{(i)} \leq Z_B^{(ij)} \leq \dots \leq Z$$

Experiments

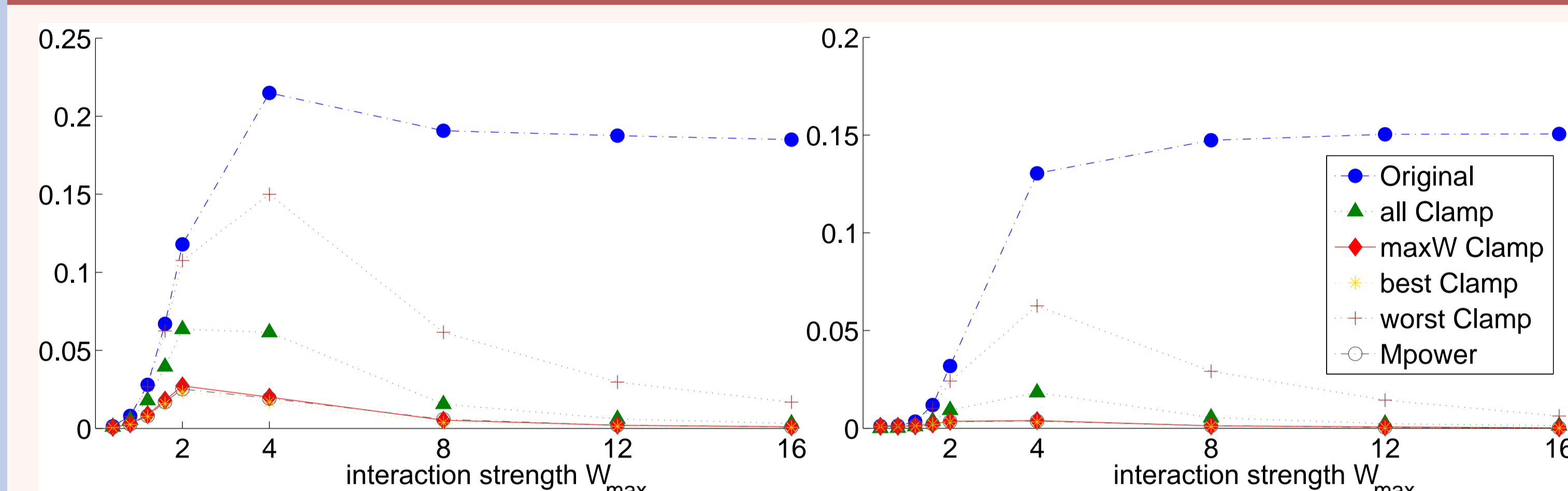
We investigate error of $Z_B^{(i)} = Z_B|_{X_i=0} + Z_B|_{X_i=1}$ compared to error of original Z_B , various methods of choosing variable to clamp X_i :

- **best Clamp** best improvement in error of Z in hindsight
- **worst Clamp** worst improvement in error of Z in hindsight
- **avg Clamp** average performance
- **maxW** max sum of incident edge weights $\sum_{j \in \mathcal{N}(i)} |W_{ij}|$
- **Mpower** more sophisticated, tries to break heavy cycles



- Potentials drawn at random: unary $\theta_i \sim U[-2, 2]$, edge $W_{ij} \sim U[-W_{max}, W_{max}]$ for *general*, $W_{ij} \sim U[0, W_{max}]$ for *attractive* models

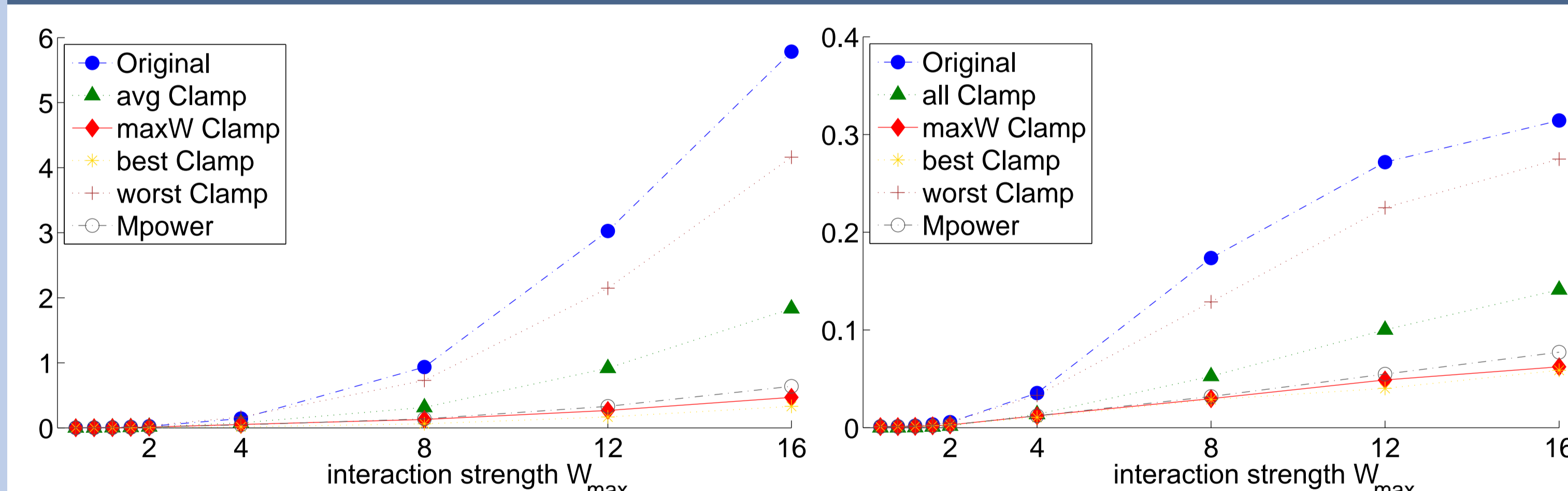
Attractive random graph $n = 10, p = 0.5$



Error of $\log Z$

Avg l_1 error marginals

General random graph $n = 10, p = 0.5$



As n grows, still helpful even just to clamp one variable (see paper)

General 'lamp' graph

